## Lecture 6: More Lists

- Theory
- Define append/3, a predicate for concatenating two lists, and illustrate what can be done with it
- Discuss two ways of reversing a list
- A naïve way using append/3
- A more efficient method using accumulators



## Lecture 6: More Lists

- Exercises
- Exercises of LPN: 6.1, 6.2, 6.3, 6.4, 6.5, 6.6
- Practical work



## Append

- We will define an important predicate append/3 whose arguments are all lists
- Declaratively, append(L1,L2,L3) is true if list L3 is the result of concatenating the lists L1 and L2 together
?- append([a,b,c,d],[3,4,5],[a,b,c,d,3,4,5]).
yes
?- append([a,b,c],[3,4,5],[a,b,c,d,3,4,5]).
no


## Append, viewed procedurally

- From a procedural perspective, the most obvious use of append/3 is to concatenate two lists together
- We can do this simply by using a variable as third argument
?- append([a,b,c,d],[1,2,3,4,5], X).
$X=[a, b, c, d, 1,2,3,4,5]$
yes
?-


## Definition of append/3

```
append([ ], L, L).
append([H|L1], L2, [H|L3]):-
    append(L1, L2, L3).
```

- Recursive definition
- Base clause: appending the empty list to any list produces that same list
- The recursive step says that when concatenating a non-empty list $[\mathrm{H} \mid \mathrm{T}]$ with a list L , the result is a list with head H and the result of concatenating T and L


## How append/3 works

- Two ways to find out:
- Use trace/0 on some examples
- Draw a search tree! Let's consider a simple example
?- append([a,b,c],[1,2,3], R).


## Search tree example

## ?- append([a,b,c],[1,2,3], R).

append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).

## Search tree example

?- append([a,b,c],[1,2,3], R). /
append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).

## Search tree example

?- append([a,b,c],[1,2,3], R).
$\dagger$
$\mathrm{R}=[\mathrm{a} \mid \mathrm{R} 0]$
?- append([b,c],[1,2,3],R0)
append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).

## Search tree example

?- append([a,b,c],[1,2,3], R).
$\dagger$

$$
\mathrm{R}=[\mathrm{a} \mid \mathrm{R} 0]
$$

?- append([b,c],[1,2,3],R0)
append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).

## Search tree example

?- append([a,b,c],[1,2,3], R).
$\dagger$
$R=[a \mid R 0]$
?- append([b,c],[1,2,3],R0)
$\dagger \quad \mathrm{R} 0=[\mathrm{b} \mid \mathrm{R} 1]$
?- append([c],[1,2,3],R1)

## Search tree example

?- append([a,b,c],[1,2,3], R).
$\dagger$

append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).
$\dagger \quad \mathrm{R} 0=[\mathrm{b} \mid \mathrm{R} 1]$
?- append([c],[1,2,3],R1)
/ ।

## Search tree example

?- append([a,b,c],[1,2,3], R).
$\dagger$

| 1 |  |
| :---: | :---: |
| $\mathrm{R}=[\mathrm{a} \mid \mathrm{R} 0]$ |  |
| ?- append([b,c],[1,2,3],R0) |  |
| 1 | , |
| $\dagger$ | $\mathrm{R} 0=[\mathrm{b} \mid \mathrm{R} 1]$ |

append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).
?- append([c],[1,2,3],R1)

$\dagger$
$\mathrm{R} 0=[\mathrm{b} \mid \mathrm{R} 1]$
$\mathrm{R} 1=[\mathrm{c} \mid \mathrm{R} 2]$
?- append([],[1,2,3],R2)

## Search tree example

?- append([a,b,c],[1,2,3], R).
$\dagger$

append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).
?- append([b,c],[1,2,3],R0)

R0=[b|R1]
?- append([c],[1,2,3],R1)

$\dagger$
$\mathrm{R} 1=[\mathrm{c} \mid \mathrm{R} 2]$
?- append([],[1,2,3],R2)
/

## Search tree example

?- append([a,b,c],[1,2,3], R).
$\dagger$

append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).
$\mathrm{R} 0=[\mathrm{b} \mid \mathrm{R} 1]$
?- append([c],[1,2,3],R1)

$\dagger$
R1=[c|R2]
?- append([],[1,2,3],R2)

R2=[1,2,3]

## Search tree example

?- append([a,b,c],[1,2,3], R).

?- append([b,c],[1,2,3],R0)
$\dagger \quad \mathrm{R} 0=[\mathrm{b} \mid \mathrm{R} 1]$
?- append([c],[1,2,3],R1)
$\dagger$
R1 $=[\mathrm{c} \mid \mathrm{R} 2]=[\mathrm{c}, 1,2,3]$
$\mathbf{R O}=[\mathrm{b} \mid \mathrm{R} 1]=[\mathrm{b}, \mathrm{c}, 1,2,3]$
$R=[a \mid R 0]=[a, b, c, 1,2,3]$
append([], L, L).
append([H|L1], L2, [H|L3]):append(L1, L2, L3).

## Using append/3

- Now that we understand how append/3 works, let's look at some applications
- Splitting up a list:

$$
\begin{array}{ll}
\text { ?- append }(X, Y,[a, b, c, d]) . \\
X=[] & Y=[a, b, c, d] ; \\
X=[a] & Y=[b, c, d] ; \\
X=[a, b] & Y=[c, d] ; \\
X=[a, b, c] & Y=[d] ; \\
X=[a, b, c, d] & Y=[] ; \\
\text { no } &
\end{array}
$$

## Prefix and suffix

- We can also use append/3 to define other useful predicates
- A nice example is finding prefixes and suffixes of a list



## Definition of prefix/2

## prefix(P,L):append( $\mathrm{P}, \mathrm{Z}, \mathrm{L})$.

- A list $P$ is a prefix of some list $L$ when there is some list such that $L$ is the result of concatenating $P$ with that list.
- We use the anonymous variable because we don't care what that list is.


## Use of prefix/2

## prefix(P,L):$\operatorname{append}(P,, L)$.

```
?- prefix(X, [a,b,c,d]).
X=[ ];
X=[a];
X=[a,b];
X=[a,b,c];
X=[a,b,c,d];
no
```


## Definition of suffix/2

$$
\begin{aligned}
& \text { suffix(S,L):- } \\
& \text { append(_,S,L). }
\end{aligned}
$$

- A list $S$ is a suffix of some list $L$ when there is some list such that $L$ is the result of concatenating that list with $S$.
- Once again, we use the anonymous variable because we couldn't care less what that list is.


## Use of suffix/2

## suffix(S,L):-

 append(_,S,L).$$
\begin{aligned}
& ?-\text { suffix }(X,[a, b, c, d]) . \\
& X=[a, b, c, d] ; \\
& X=[b, c, d] ; \\
& X=[c, d] ; \\
& X=[d] ; \\
& X=[] ; \\
& \text { no }
\end{aligned}
$$

## Definition of sublist/2

- Now it is very easy to write a predicate that finds sub-lists of lists
- The sub-lists of a list $L$ are simply the prefixes of suffixes of $L$
sublist(Sub,List):suffix(Suffix,List), prefix(Sub,Suffix).


## append/3 and efficiency

- The append/3 predicate is useful, and it is important to know how to use it
- It is of equal importance to know that append/3 can be source of inefficiency
- Why?
- Concatenating a list is not done in one simple action
- But by traversing down one of the lists


## Question

- Using append/3 we would like to concatenate two lists:
- List 1: [a,b,c,d,e,f,g,h,i]
- List 2: [j,k,l]
- The result should be a list with all the elements of list 1 and 2, the order of the elements is not important
- Which of the following goals is the most efficient way to concatenate the lists?
?- append([a,b,c,d,e,f,g,h,i],[j,k,l],R).
?- append([j,k,l],[a,b,c,d,e,f,g,h,i],R).


## Answer

- Look at the way append/3 is defined
- It recurses on the first argument, not really touching the second argument
- That means it is best to call it with the shortest list as first argument
- Of course you don't always know what the shortest list is, and you can only do this when you don't care about the order of the elements in the concatenated list
- But if you do, it can help make your Prolog code more efficient


## Reversing a List

- We will illustrate the problem with append/3 by using it to reverse the elements of a list
- That is, we will define a predicate that changes a list [a,b,c,d,e] into a list [e,d,c,b,a]
- This would be a useful tool to have, as Prolog only gives easy access to the front of the list


## Naïve reverse

## Recursive definition

1. If we reverse the empty list, we obtain the empty list
2. If we reverse the list [H|T], we end up with the list obtained by reversing T and concatenating it with $[\mathrm{H}]$
To see that this definition is correct, consider the list [a,b,c,d].

- If we reverse the tail of this list we get [d,c,b].
- Concatenating this with [a] yields [d,c,b,a]


## Naïve reverse in Prolog

## naiveReverse([],[]).

naiveReverse([H|T],R):-
naiveReverse(T,RT),
append(RT,[H],R).

- This definition is correct, but it does an awful lot of work
- It spends a lot of time carrying out appends
- But there is a better way...


## Reverse using an accumulator

- The better way is using an accumulator



## Reverse using an accumulator

- The accumulator will be a list, and when we start reversing it will be empty

$\square$ Me continue this until we reach the empty list At this point the accumulator will contain the reversed list!


## Reverse using an accumulator

- We simply take the head of the list that we want to reverse and add it to the head of the accumulator list


## Reverse using an accumulator



- We continue this until we reach the empty list
 contain the reversed list!


## Reverse using an accumulator



## Reverse using an accumulator

accReverse([ ],L,L). accReverse([H|T],Acc,Rev):accReverse(T,[H|Acc],Rev).

## Adding a wrapper predicate

accReverse([ ],L,L). accReverse([H|T],Acc,Rev):accReverse(T,[H|Acc],Rev).
reverse(L1,L2):-
accReverse(L1,[ ],L2).

## Illustration of the accumulator

- List: [a,b,c,d] Accumulator: []


## Illustration of the accumulator

- List: [a,b,c,d] Accumulator: []
- List: [b,c,d] Accumulator: [a]


## Illustration of the accumulator

- List: [a,b,c,d] Accumulator: []
- List: [b,c,d] Accumulator: [a]
- List: [c,d] Accumulator: [b,a]


## Illustration of the accumulator

- List: [a,b,c,d] Accumulator: []
- List: [b,c,d] Accumulator: [a]
- List: [c,d] Accumulator: [b,a]
- List: [d]

Accumulator: [c,b,a]

## Illustration of the accumulator

- List: [a,b,c,d] Accumulator: []
- List: [b,c,d] Accumulator: [a]
- List: [c,d] Accumulator: [b,a]
- List: [d]
- List: []

Accumulator: [c,b,a]
Accumulator: [d,c,b,a]

## Summary of this lecture

- The append/3 is a useful predicate, don't be scared of using it
- However, it can be a source of inefficiency
- The use of accumulators is often better
- We will encounter a very efficient way of concatenating list in later lectures, where we will explore the use of "difference lists"


## Next lecture

- Definite Clause Grammars
- Introduce context free grammars and some related concepts
- Introduce DCGs, definite clause grammars, a built-in Prolog mechanism for working with context free grammars

